

Lecture

Minimum Spanning trees

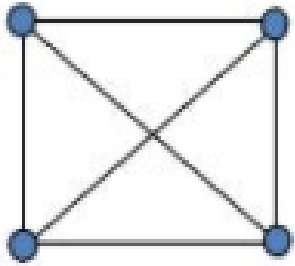
Prim's algorithm

Spanning Trees

A connected sub graph H of given Graph G is said to be spanning tree iff

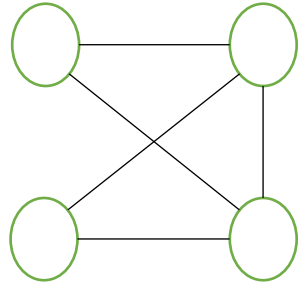
1. H contains all vertices of G.
2. H should contain $(n-1)$ edges if G contains n vertices.

Question 1 Find number of spanning trees for the following graph:



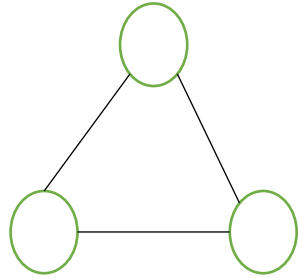
Spanning Trees

Question 2 Find number of spanning trees for the following graph:



Spanning Trees

Question 3 Find number of spanning trees for the following graph:



Spanning Trees

The number of spanning trees in a complete graph:

$$ST(K_n) = n^{n-2}$$

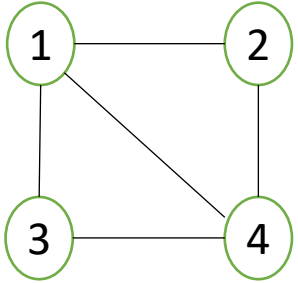
$$ST(K_2) =$$

$$ST(K_3) =$$

$$ST(K_4) =$$

$$ST(K_5) =$$

Spanning Trees: Kirchoff's Theorem



Find number of spanning trees for the following graph:

Step 1: Find adjacency matrix for the given graph.

$$M = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

Step 2: Replace all diagonal 0's by its' degree and non-diagonal 1's by -1.

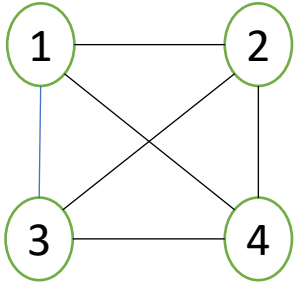
$$M = \begin{bmatrix} 3 & -1 & -1 & -1 \\ -1 & 2 & 0 & -1 \\ -1 & 0 & 2 & -1 \\ -1 & -1 & -1 & 3 \end{bmatrix}$$

Step 3: Number of spanning tree = co-factor of any element.

$$\text{Co-factor } (M_{11}) = \begin{vmatrix} 2 & 0 & -1 \\ 0 & 2 & -1 \\ -1 & -1 & 3 \end{vmatrix}$$

$$= 2(6-1)-1(2) = 8$$

Spanning Trees: Kirchoff's Theorem



Find number of spanning trees for the following graph:

Step 1: Find adjacency matrix for the given graph.

$$M = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

Step 2: Replace all diagonal 0's by its' degree and non-diagonal 1's by -1.

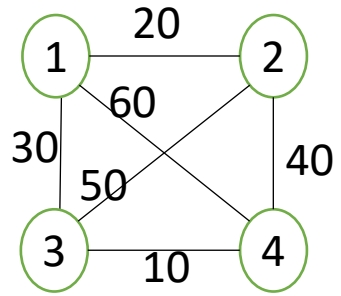
$$M = \begin{bmatrix} 3 & -1 & -1 & -1 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ -1 & -1 & -1 & 3 \end{bmatrix}$$

Step 3: Number of spanning tree = co-factor of any element.

$$\text{Co-factor } (M_{11}) = \begin{vmatrix} 3 & -1 & -1 \\ -1 & 3 & -1 \\ -1 & -1 & 3 \end{vmatrix}$$

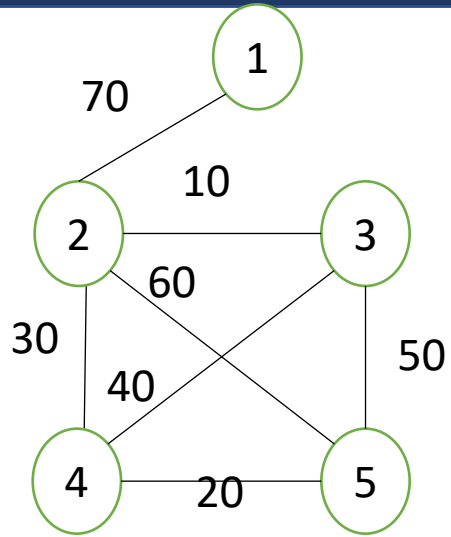
$$= 3(9-1) + 1(-3-1) - 1(1+3) = 16$$

Minimum Cost Spanning Tree



Using Prim's and Kruskal's algorithm we can find out the minimum cost spanning tree

Minimum Cost Spanning Tree: Prim's Algorithm

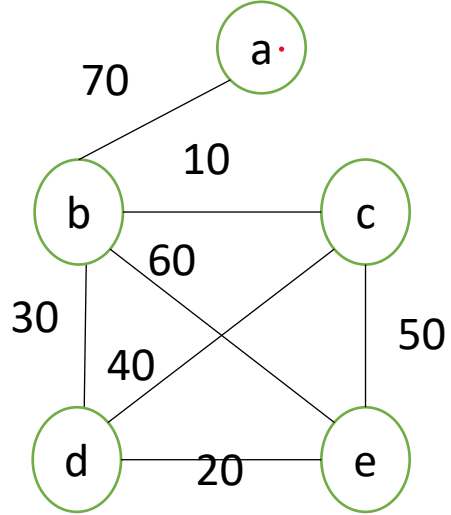


Step 1: Choose any vertex and find adjacent of that vertex and take minimum of them

Step 2: Find adjacency vertices of new vertex and take minimum of these adjacent and previous.

Step 3: Repeat step 2 until all the vertices are not covered.

Minimum Cost Spanning Tree: Prim's Algorithm



The edges in set A always form a single tree

Starts from an arbitrary “root”: $V_A = \{a\}$

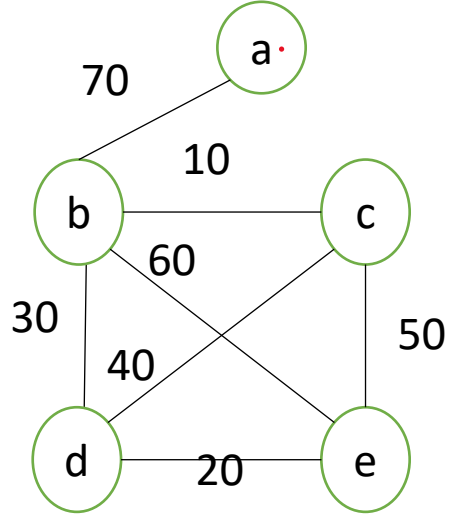
At each step:

Find a light edge crossing $(V_A, V - V_A)$

Add this edge to A

Repeat until the tree spans all vertices

Minimum Cost Spanning Tree: Prim's Algorithm



The edges in set A always form a single tree

Starts from an arbitrary “root”: $V_A = \{a\}$

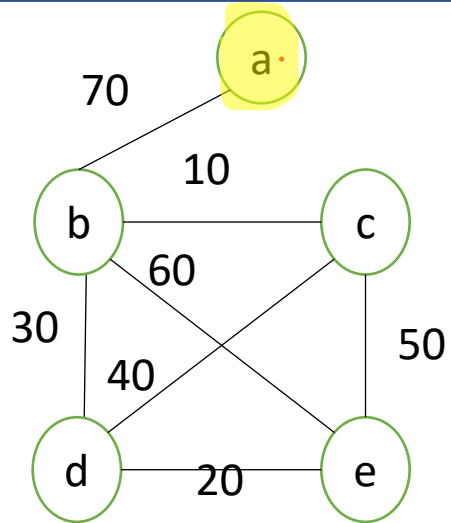
At each step:

Find a light edge crossing $(V_A, V - V_A)$

Add this edge to A

Repeat until the tree spans all vertices

Minimum Cost Spanning Tree: Prim's Algorithm



Nodes distances in priority Queue : 0 ∞ ∞ ∞ ∞

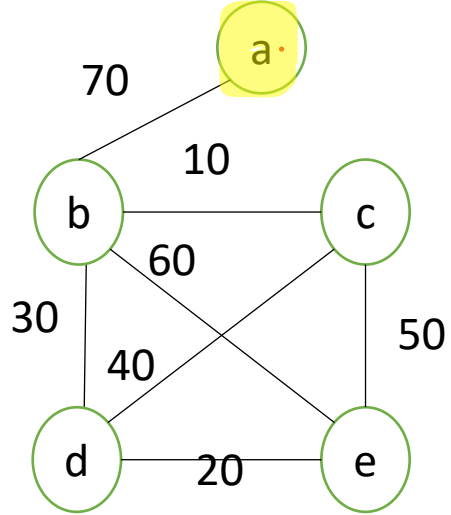
$Q = \{a, b, c, d, e\}$

$V_A = \emptyset$

Extract-MIN(Q) \Rightarrow a

$V_A = \{a\}$

Minimum Cost Spanning Tree: Prim's Algorithm



$\text{key}[b] = 70 \quad \pi[b] = a$

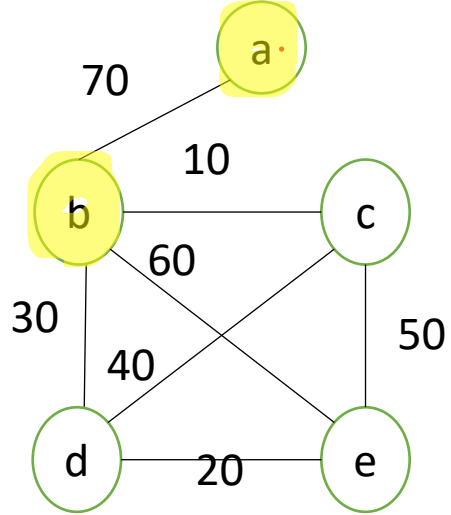
70 $\infty \infty \infty \infty$

$Q = \{b, c, d, e\} \quad V_A = \{a\}$

$\text{Extract-MIN}(Q) \Rightarrow b$

$V_A = \{a, b\}$

Minimum Cost Spanning Tree: Prim's Algorithm



$\text{key}[b] = 70 \quad \pi[b] = a$

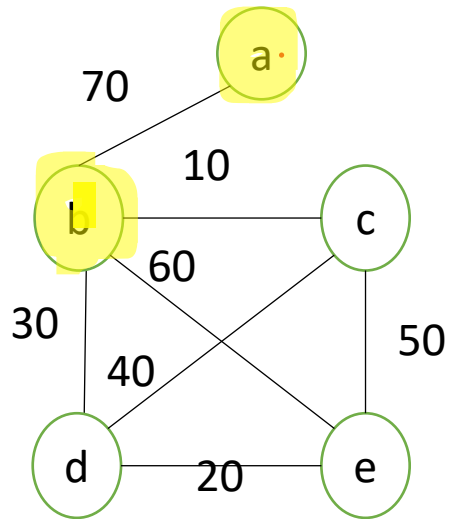
70 $\infty \infty \infty \infty$

$Q = \{b, c, d, e\} \quad V_A = \{a\}$

$\text{Extract-MIN}(Q) \Rightarrow b$

$V_A = \{a, b\}$

Minimum Cost Spanning Tree: Prim's Algorithm



key [c] = 10 π [c] = b

key [d] = 30 π [d] = b

key [e] = 60 π [e] = b

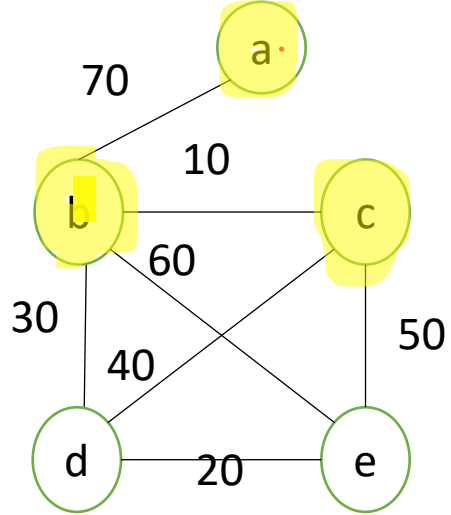
10, 30, 60

$Q = \{c, d, e\}$ $V_A = \{a, b\}$

Extract-MIN(Q) \Rightarrow c

$V_A = \{a, b, c\}$

Minimum Cost Spanning Tree: Prim's Algorithm



key [c] = 10 π [c] = b

key [d] = 30 π [d] = b

key [e] = 60 π [e] = b

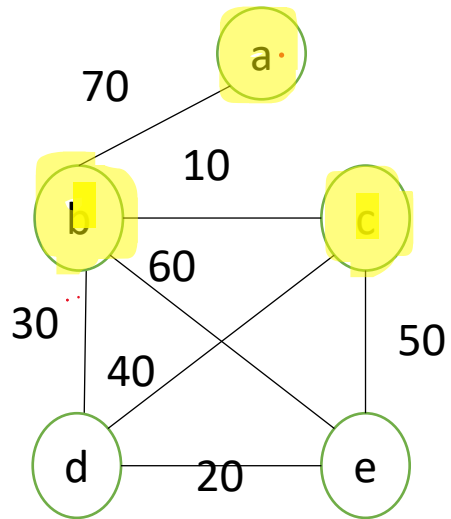
10, 30, 60

$Q = \{c, d, e\}$ $V_A = \{a, b\}$

Extract-MIN(Q) \Rightarrow c

$V_A = \{a, b, c\}$

Minimum Cost Spanning Tree: Prim's Algorithm



key [d] = 30. π [d] = b

key [e] = 50 π [e] = c

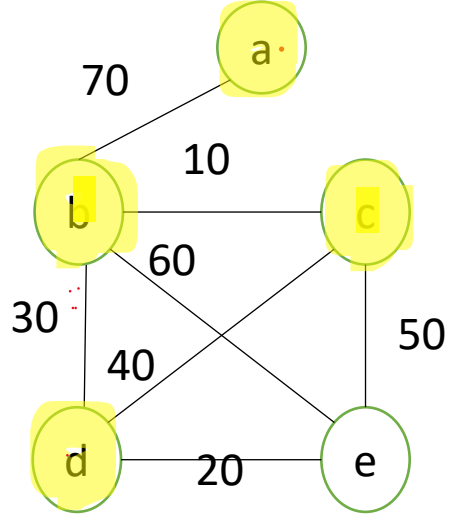
30, 50

$Q = \{d, e\}$ $V_A = \{a, b, c\}$

Extract-MIN(Q) \Rightarrow d

$V_A = \{a, b, c, d\}$

Minimum Cost Spanning Tree: Prim's Algorithm



key [d] = 30. π [d] = b

key [e] = 50 π [e] = c

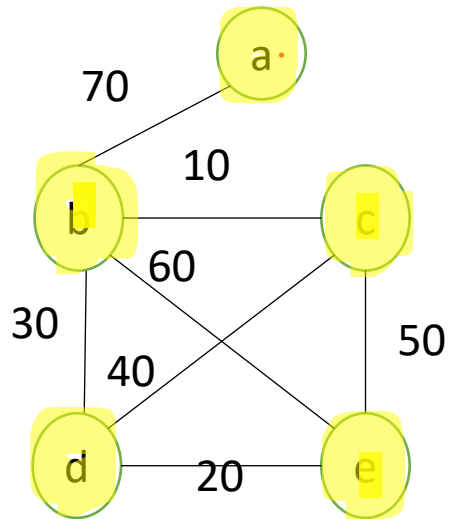
30, 50

$Q = \{d, e\}$ $V_A = \{a, b, c\}$

Extract-MIN(Q) \Rightarrow d

$V_A = \{a, b, c, d\}$

Minimum Cost Spanning Tree: Prim's Algorithm



key [e] = 20 π [e] = d

20

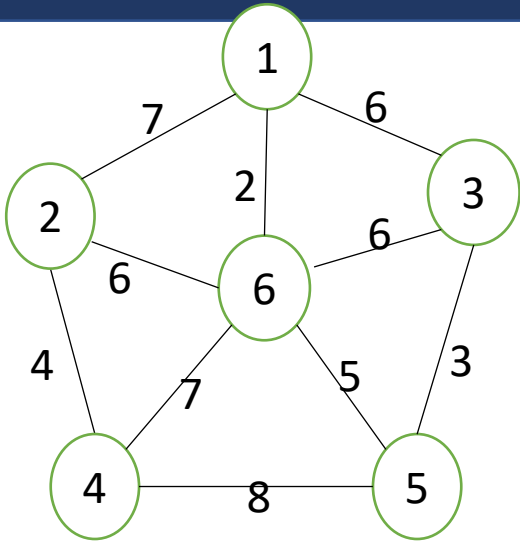
$Q = \{ e \}$ $V_A = \{ a, b, c, d \}$

Extract-MIN(Q) \Rightarrow e

$V_A = \{ a, b, c, d, e \}$

$Q = \emptyset$

Minimum Cost Spanning Tree: Prim's Algorithm



Minimum Cost Spanning Tree: Prim's Algorithm

MST-PRIM(G, w, r)

```
1  for each  $u \in G.V$ 
2       $u.key = \infty$ 
3       $u.\pi = \text{NIL}$ 
4   $r.key = 0$ 
5   $Q = G.V$ 
6  while  $Q \neq \emptyset$ 
7       $u = \text{EXTRACT-MIN}(Q)$ 
8      for each  $v \in G.Adj[u]$ 
9          if  $v \in Q$  and  $w(u, v) < v.key$ 
10              $v.\pi = u$ 
11              $v.key = w(u, v)$ 
```

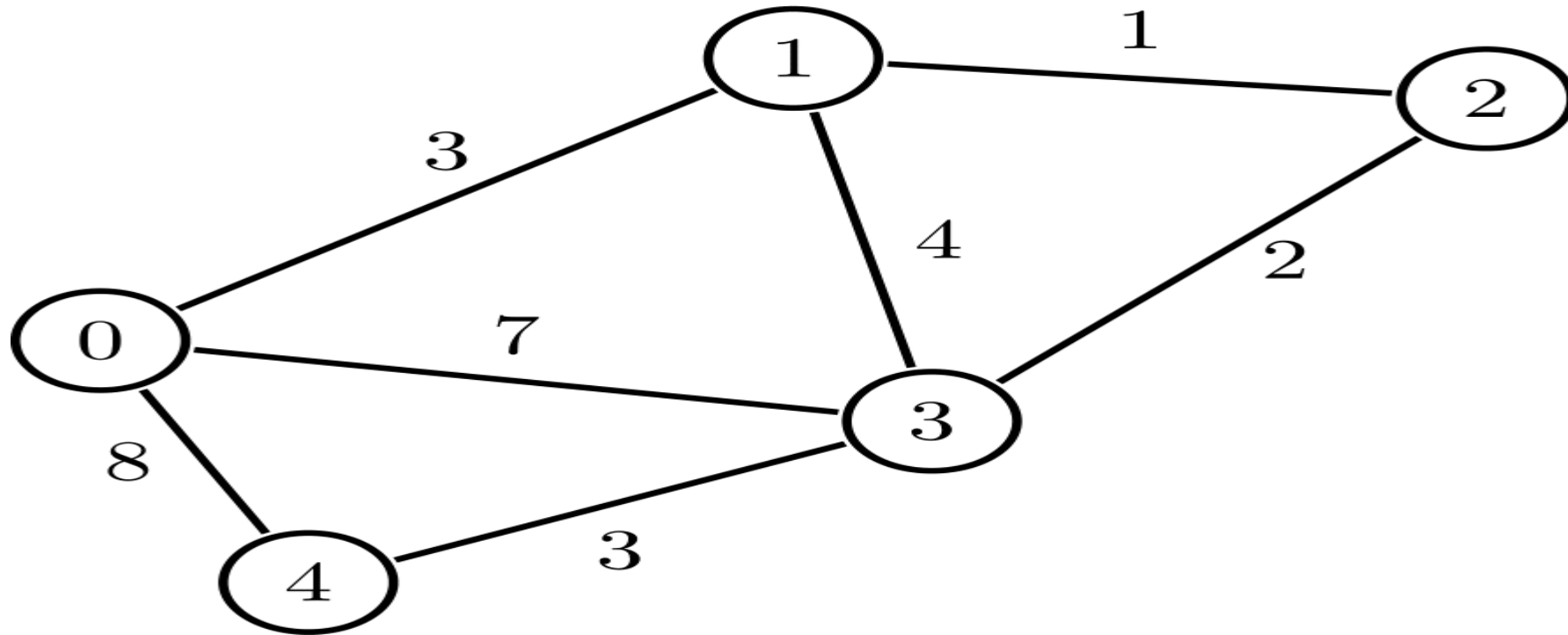
Time Complexity:

Using Binary Min Heap: $O[(V + E)\log V]$

Using Fibonacci Min Heap: $O[E + V\log V]$

Using Binomial Min Heap: $O[(V + E)]$

Minimum Cost Spanning Tree: Prim's Algorithm





Thank You