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# OPERATING SYSTEM: CSET209



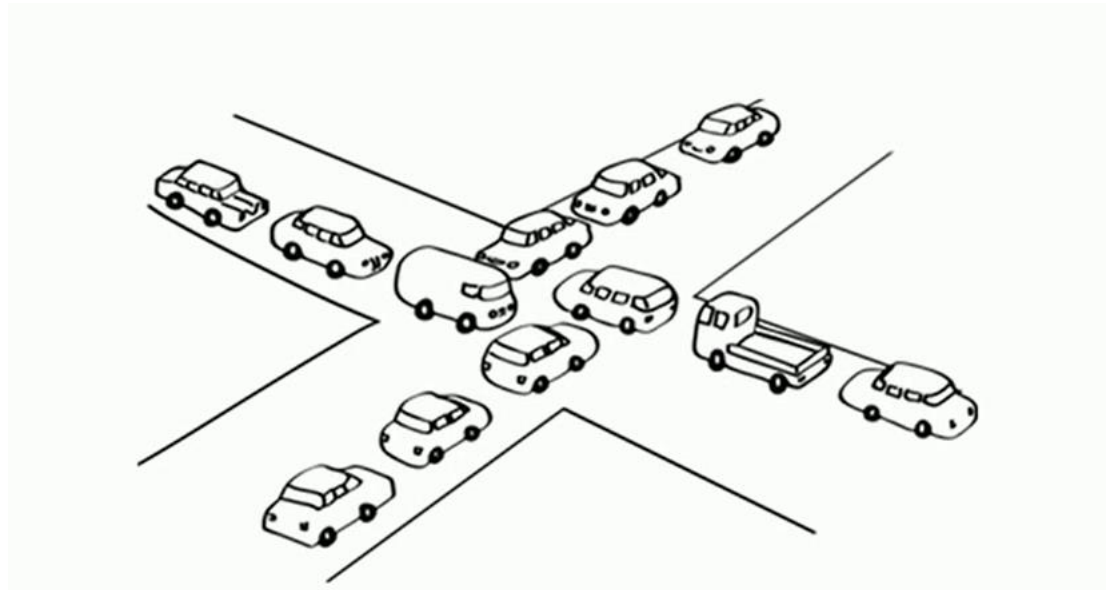
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# OUTLINE

- Deadlock
- Resource Allocation graphs

# DEADLOCK

A situation where a set of processes **continue to run indefinitely** without making any **progress**. In this situation, none of the process gets executed since the **resource it needs, is held by some other process** which is also **waiting for some other resource to be released**.



If a deadlock occurs, it can be resolved if one car backs up (preempt resources and rollback)

# DEADLOCK EXAMPLE

- A set of blocked processes each holding a resource and waiting to acquire a resource held by another process in the set
- Example
  - semaphores  $A$  and  $B$ , initialized to 1

$P_0$

wait (A);

wait (B);

$P_1$

wait(B)

wait(A)

# DEADLOCK EXAMPLE

```
void transaction(Account from, Account to, double amount)
```

```
{  
    mutex lock1, lock2;  
    lock1 = get_lock(from);  
    lock2 = get_lock(to);  
    acquire(lock1);  
        acquire(lock2);  
            withdraw(from, amount);  
            deposit(to, amount);  
        release(lock2);  
    release(lock1);  
}
```

- Transactions 1 and 2 execute concurrently. Transaction 1 transfers \$25 from account A to account B, and Transaction 2 transfers \$50 from account B to account A

# NECESSARY CONDITIONS FOR DEADLOCKS

## **1. Mutual Exclusion**

A resource can only be shared in mutually exclusive manner. It implies, if two process cannot use the same resource at the same time.

## **2. Hold and Wait**

A process waits for some resources while holding another resource at the same time.

## **3. No preemption**

A resource can be released only voluntarily by the process holding it, after that process has completed its task

## **4. Circular Wait**

All the processes must be waiting for the resources in a cyclic manner so that the last process is waiting for the resource which is being held by the first process.

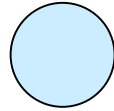
# RESOURCE-ALLOCATION GRAPH

A directed graph used to model resource requests and allocations.

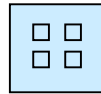
**request edge** – directed edge  $P_i \rightarrow R_j$

**assignment edge** – directed edge  $R_j \rightarrow P_i$

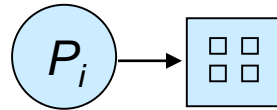
Process



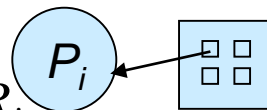
Resource Type with 4 instances



$P_i$  requests instance of  $R_j$

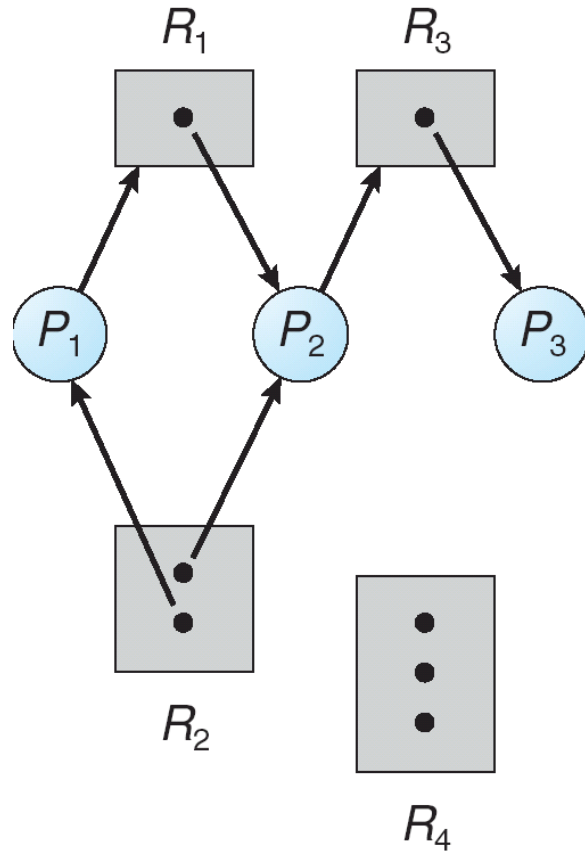


$P_i$  is holding an instance of  $R_j$



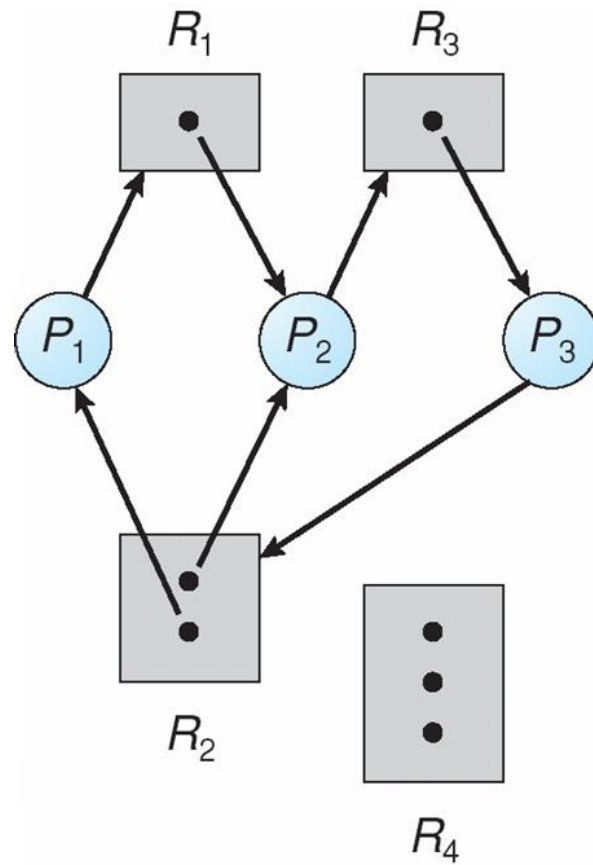
$R_j$

# RESOURCE-ALLOCATION GRAPH EXAMPLE



- **IS THERE A DEADLOCK HERE?**
- **NO**

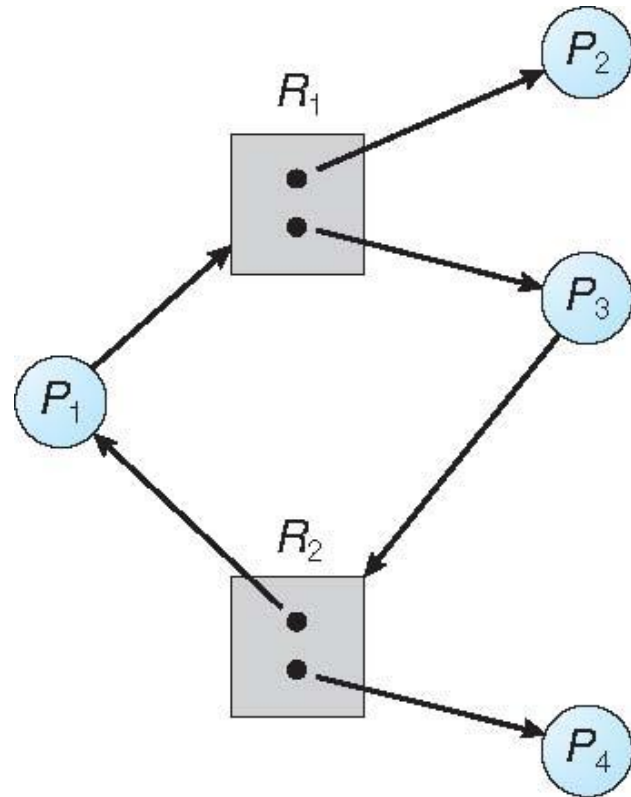
# RESOURCE-ALLOCATION GRAPH EXAMPLE



IS THERE A DEADLOCK HERE?

YES

# RESOURCE-ALLOCATION GRAPH EXAMPLE



**IS THERE A DEADLOCK HERE?**

**NO**

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## BASIC FACTS

- If graph contains no cycles  $\Rightarrow$  no deadlock
- If graph contains a cycle  $\Rightarrow$ 
  - if only one instance per resource type, then deadlock
  - if several instances per resource type, possibility of deadlock

# DEADLOCK IS A PROBABILISTIC EVENT

□ Suppose there are 3 processes A, B, and C requesting resources R, S, and T as follows.

A	B	C
Request R	Request S	Request T
Request S	Request T	Request R
Release R	Release S	Release T
Release S	Release T	Release R

□ Let the order of execution be

A request R

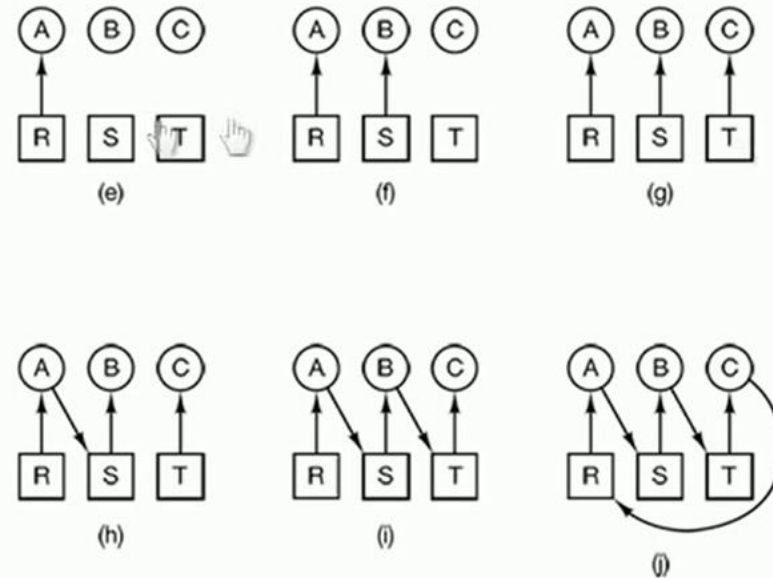
B request S

C request T

A request S

B request T

C request R



(deadlock)

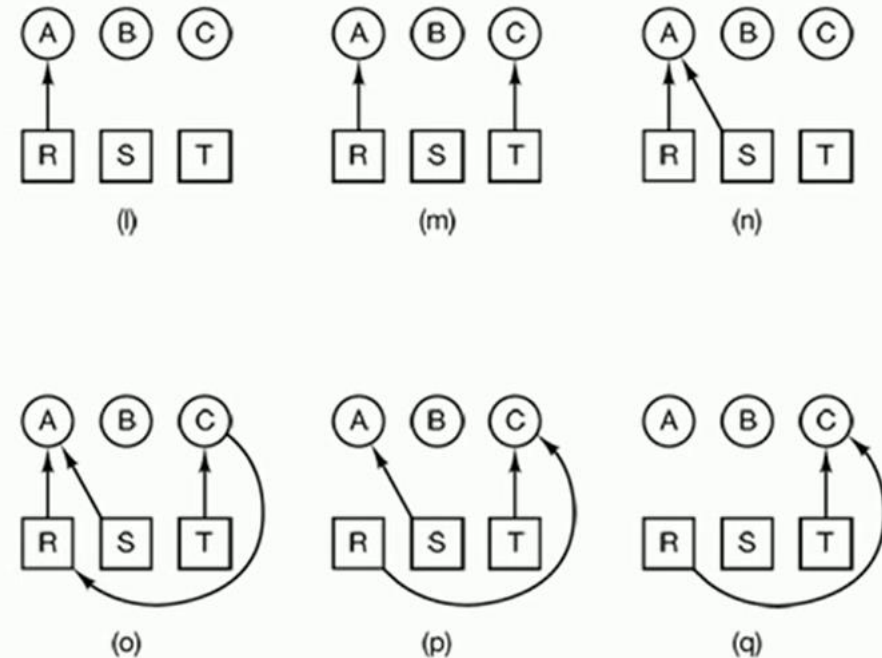
# DEADLOCK IS A PROBABILISTIC EVENT

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Request R	Request S	Request T
Request S	Request T	Request R
Release R	Release S	Release T
Release S	Release T	Release R

□ Let the order of execution be

- A request R
- C request T
- A request S
- C request R
- A release R
- A release S



NO deadlock)



Thank you !